Performance evaluation of the TIANC protocol



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2nd December 2015

Based on *Inducing Collisions for Fast RFID Tag Identification* by Gianluca Carroccia and Gaia Maselli

Introduction



Recap: RFID systems

- Goal: identification of RFID tags
- One reader queries many tags
- Tags are passive: powered by the reader, no collision detection
- Reader arbitrates the channel using some MAC protocol
- Traditional protocols try to minimize collisions
 - Anti-collision protocols

Recap: Analog Network Coding (ANC)

- Since we can't really avoid collisions, let's exploit them!
- By using ANC we can recover the original transmission, even if there was a collision
- Basic assumption:
 - Reader has multiple antennas
 - Each antenna receives a signal with some attenuation coefficients
 - Reader can create and solve a system of equations, by combining all received signals.

Description of the TIANC protocol



TIANC: Tag Identification through ANC

- Reader has a number a > 1 of antennas:
 - Only a single antenna queries the tags
 - All the antennas receive tags tranmissions
 - Their interrogation zones overlap



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TIANC: Tag Identification through ANC

- Reader has a number *a* > 1 of antennas:
 - Only a single antenna queries the tags
 - All the antennas receive tags tranmissions
 - Their interrogation zones overlap
- Number *n* of tags is supposed to be known
- If there are $k \leq a$ collisions:

$$\begin{cases} a_1 = b_{1,1} \cdot s_1 + b_{1,2} \cdot s_2 + \dots + b_{1,k} \cdot s_k \\ \vdots \\ a_a = b_{a,1} \cdot s_1 + b_{a,2} \cdot s_2 + \dots + b_{a,k} \cdot s_k \end{cases}$$

 The reader can recover the signals s₁,..., s_k and demodulate them into the k tags ID.

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TIANC frames

- TIANC is based on Tree Slotted Aloha (TSA)
- Each frame is divided into Training and Identification subframes
- Tags transmit in both subframes



Training subframe

Two goals:

- Goal 1: estimate channel coefficients
 - Tags send one bit with content 1 in a random slot
 - Reader can estimate and store the attenuation coefficients



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- Goal 2: group tags into colliding sets



Training subframe

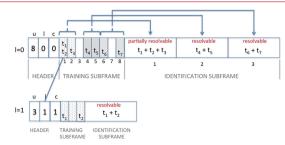
Two goals:

- Goal 1: estimate channel coefficients
 - Tags send one bit with content 1 in a random slot
 - Reader can estimate and store the attenuation coefficients
- Goal 2: group tags into colliding sets
- At the end of each slot, the reader replies with another bit:
 - ACK = singleton or collision slot (reader cannot distinguish!)
 - \circ NACK = idle slot

Identification subframe

- Tags in the same group transmit in the same identification slot
- There are only colliding slots! Either *resolvable* or *unresolvable*
- Resolvable slots:
 - Exactly a tags transmit
 - Reader can solve a system of a equations and recover all the a IDs
- Unresolvable (or partially resolvable) slots
 - More than a tags transmit
 - Reader cannot decode all the ID
- For each unresolvable slot, a new child frame has to be issued by the reader

Example of execution (a = 2)



- Idea: map a = 2 consecutive non-idle slots in the training subframe into the same slot of the identification subframe
- Mapping from slot *i* to slot *j*:

$$j = \left\lceil \frac{i - s_{idle}}{a} \right\rceil$$

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Performance Evaluation



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- Service = retrieve the 96 bit ID of all the tags within the reader interrogation zone
- Possible outcomes:
 - Zero ID retrieved, if there are no tags
 - All IDs retrieved otherwise (assuming no hardware faults)

System-specific performance metrics

- Latency: How much time it takes to identify all tags?
 A.k.a. responsiveness or response time
- System Efficiency: How much time is spent in identification
 - Not covered by the article



General Metrics

- Speed metrics:
 - Responsiveness or response time \Rightarrow Latency
 - Throughput: not covered
 - Utilization: not really the case, as there is only one service
- Reliability metrics
 - The system is not supposed to execute the service incorrectly
- Availabilty metrics:
 - The system is not supposed to be unavailable

- Time = duration of training and identification slots
- EPCglobal standard: to transmit one bit it takes:
 - 25 μs (40 kbps bit rate) + physical overhead ($T_{phy} = 275.06 \ \mu s$)
- In general, to compute the time of a slot s:

$$T_s = T_{phy} + T(\# bits in s)$$

Example: trainings slots, only 2 bits:

$$T_{tr} = T_{phy} + T_{2bits} = 325.06 \ \mu s$$



How many slots?

• Goal: we want to compute the time to identify *n* tags:

$$T(n) = H(n) \cdot T_H + X_{tr}(n) \cdot T_{tr} + X_{id}(n) \cdot T_{id}$$

- where:
 - H(n) = # frame headers
 - $X_{tr}(n) = \#$ training slots
 - $X_{id}(n) = \#$ identification slots
- Analytic method is used:
 - Due to e.g. time and cost criterions
 - Simplifying assumptions, e.g. error-free channel estimation

- Let u = α · n be the number of slots in a frame (α > 0)
- Tags choose slots randomly and independently
- Probabability of collision given by binomial distribution:

$$p(u, n, k) = \binom{n}{k} \left(\frac{1}{u}\right)^k \left(1 - \frac{1}{u}\right)^{n-k}$$

 Reader recursively allocates a new frame for each collision of k ≥ 2 tags:

$$X_{tr}(n) = \begin{cases} 1 & \text{if } n = 1\\ u + \sum_{k=2}^{n} u \cdot p(u, n, k) \cdot X_{tr}(k) & \text{if } n > 1 \end{cases}$$

- $X_{tr}(k)$ is the # of training slots in a child subframe
- How many children? One for each collision!
- $u \cdot p(u, n, k)$ is the expected # of slots when k tags collide

How many frame headers?

• Each frame requires an header:

$$H(n) = \begin{cases} 1 & \text{if } n = 1\\ 1 + \sum_{k=2}^{n} u \cdot p(u, n, k) \cdot H(k) & \text{if } n > 1 \end{cases}$$

 Size of identification subframe depends on the outcome of the previous training subframe:

$$X_{tr}(n) = \begin{cases} 1 & \text{if } n = 1\\ v(u) + \sum_{k=2}^{n} u \cdot p(u, n, k) \cdot X_{id}(k) & \text{if } n > 1 \end{cases}$$

• $u \cdot p(u, n, 0)$ is the # of idle slots

• $v(u) = \frac{u - u \cdot p(u,n,0)}{a}$ is the # of identification slots in the first subframe

- Workload = A set of *n* tags to be identified
- Workload is synthetic
- The value of *n* does not impact on protocol performance



- Depends on the outcome of training slots
- The more collisions in training subframes, the more child frames
- The bigger the training subframe, the fewer the collisions



Goal:

Find the optimal value of α such that the training frame size $u = \alpha \cdot n$ leads to the best latency



Results of the analysis (2)

$u = \alpha \cdot n$	а
2.3 <i>n</i>	2
2.1 <i>n</i>	3
2.0 <i>n</i>	4
2.0 <i>n</i>	5

Table 1: Optimal training frame size u for each number of antennas

Results of the analysis (3)

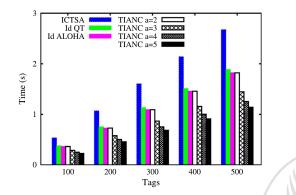


Figure 1: Comparison with ICTSA and ideal protocols (1 slot/query per tag)

Results of the analysis (4)

Speedup	а
31%	2
46%	3
53%	4
57%	5

Table 2: Speedup over ICTSA protocol



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What about Time System Efficiency?

• Recap:
$$TSE = \frac{n}{X(n) - Y(n) + \beta \cdot Y(n)}$$

•
$$X(n) = #$$
 of slots in TSA tree

•
$$Y(n) = \#$$
 of idle slots in TSA tree

 $\circ\;$ idle slots last a β fraction of non-idle slots

TSE should be defined differently for training and identification slots

- Identification slots cannot be idle!
- With training slots:
 - Similar to the TSA one, but...
 - β not so small, since tags transmit only 1 bit!



Conclusions

- TIANC features innovative elements such as its training mechanism
- Performance analysis shows significant speed improvements over previous work
- Performance not affected by large number of tags



Thanks for your attention! Questions?



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